

Dealing with Execution Uncertainty in the Continuous Double Auction

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- 1 Introduction
- 2 Mechanism
- 3 Evaluation
- 4 Discussion

What is Resource Allocation?

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To assign available resources in an economic way. [Wikipedia]

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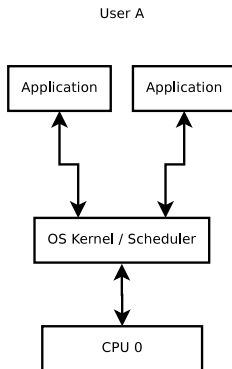
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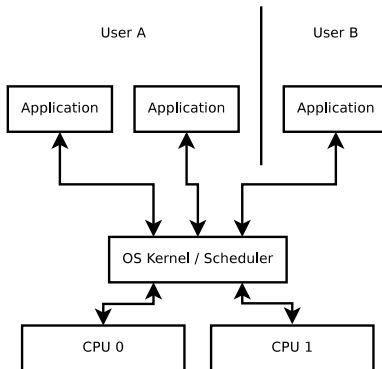
Computational Resource Allocation

- CPU time slice allocation
- Wireless channel allocation
- Network bandwidth allocation

CPU time slice allocation



CPU time slice allocation



CPU time slice allocation

This case is easy:

- There is a clear central authority (the kernel)
- The users are usually the owners of the system

Why Grid Computing?

- Processing enormous quantities of data
- Collaboratively tackling difficult challenges
- In general: computational resource sharing

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LHC Computing Grid (LCG)

- Processes 300 MB/s filtered data (from 300 GB/s raw)
- Combines 140 computing centers in 33 countries

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- Computing power available on-demand: utility computing
- Analogy: the electricity grid

Grid Resource Allocation

Grid resource allocation is harder:

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- No clear central authority

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- Decentralised resource allocation

Grid resource allocation is 'easy':

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Grid resource allocation is 'easy':

- VOs created in cooperative setting
- Hence, conflict of interest not a problem

Volunteer Computing

- Volunteers sign up, install software on their computer
- Runs when computer is idle
- Fetches and completes work packages
- Volunteer earns reputation 'points'

Why Volunteer Computing?

IBM Roadrunner

- #1 super computer
- 1.1 PetaFLOPS (130.000 cores)

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Folding@Home (Stanford)

- Protein folding research (medical)
- Utilises CPU, GPU and PS3
- 5.0 PetaFLOPS (425.000 clients)

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- So: most important tasks for project assigned first

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- Application to time-critical problems difficult

Large-scale, open, distributed systems

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Large-scale, open, distributed systems

- Combine resources in ad-hoc fashion
- May include:
 - University computing centres
 - Competing corporations
 - Consumer PCs
 - Network of workstations that are idle overnight

Problems

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- Market-based mechanisms

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- (Some) market-based mechanisms

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Continuous Double Auction

Decentralised, continuous, market-based resource allocation

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- (Some) market-based mechanisms
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Continuous Double Auction

Decentralised, continuous, market-based resource allocation

- Has the first four properties required
- Lacks trust-based decision making

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The Continuous Double Auction: Setting

- Market for trading a specific commodity
- Multiple buyers and sellers
- Buyers and sellers have their own valuations (limit prices)
- Traders shout prices (offers)
 - Bid: offer to buy
 - Ask: offer to sell

The Continuous Double Auction: Mechanism

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- Order Books

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The Continuous Double Auction: Mechanism

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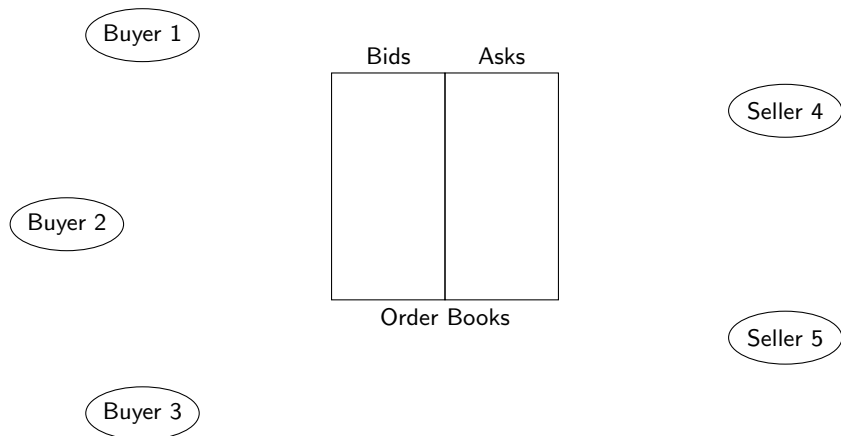
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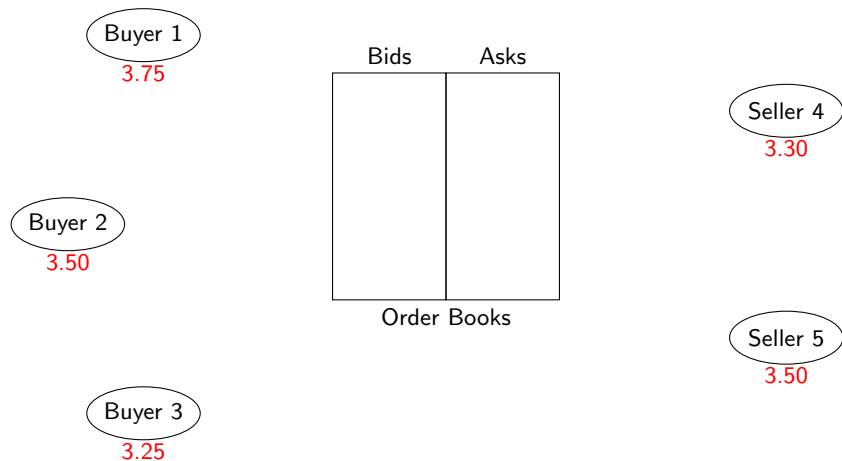
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- Clearing Rule
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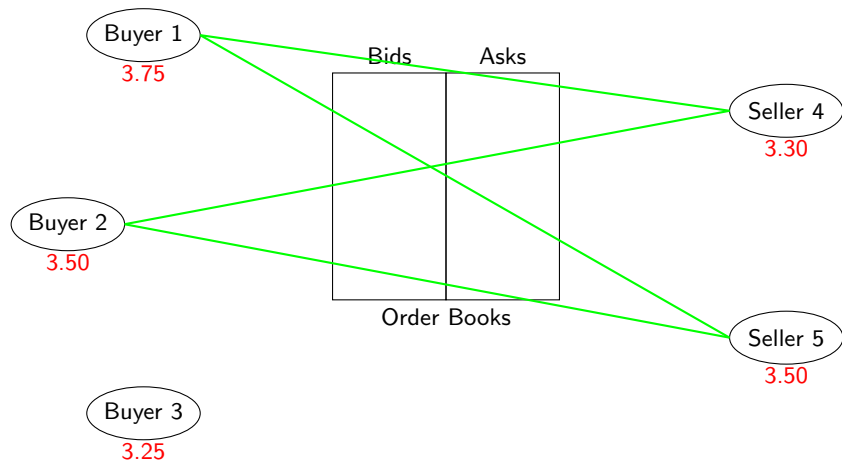
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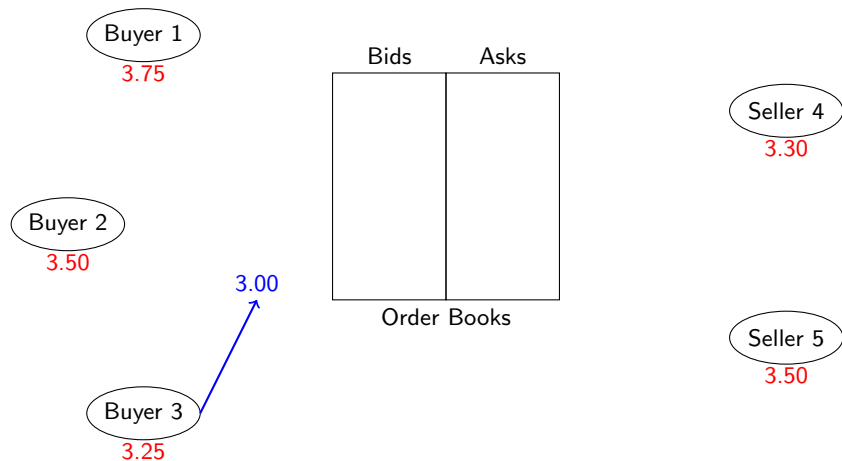
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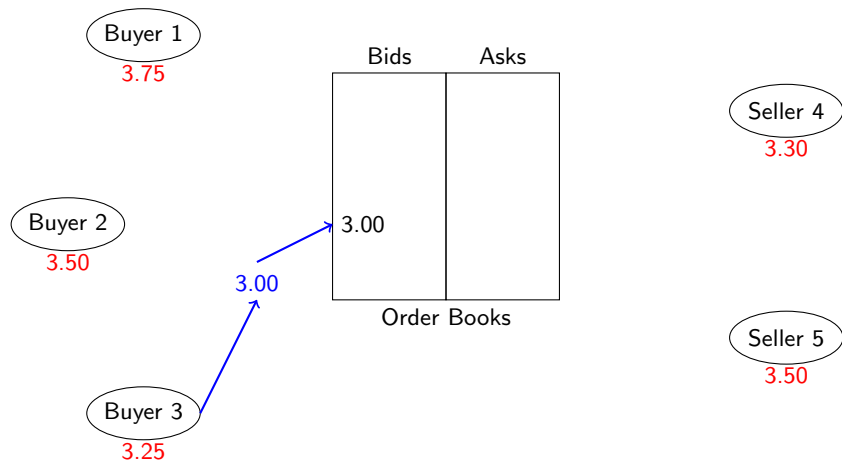
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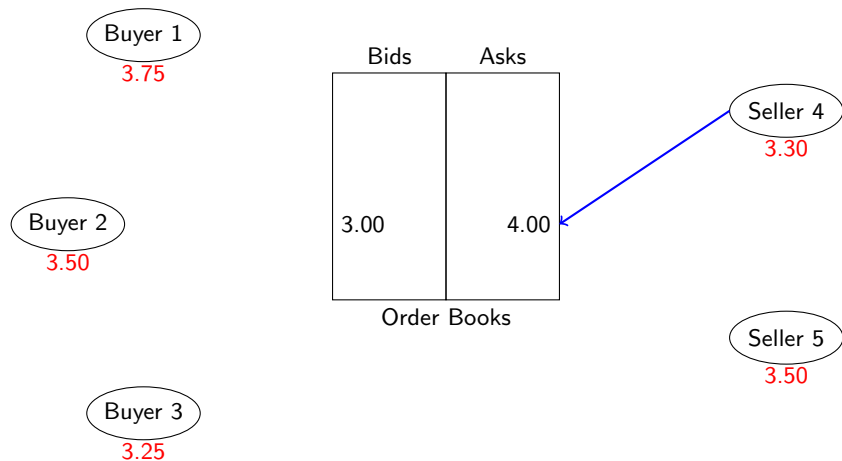
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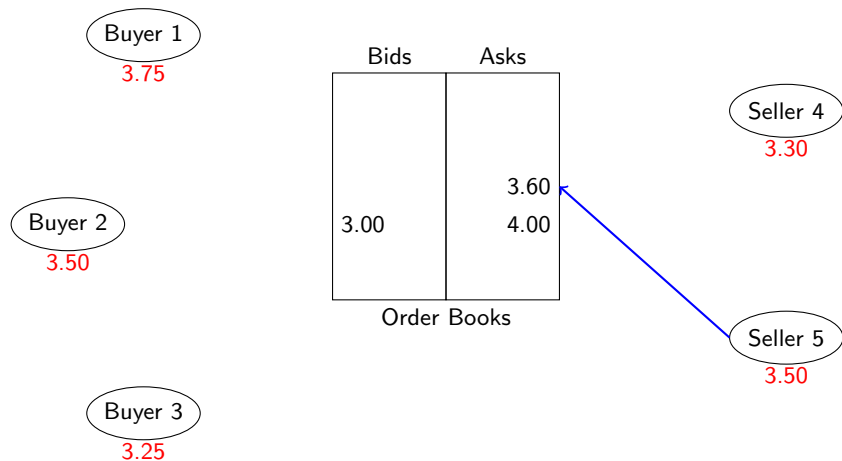
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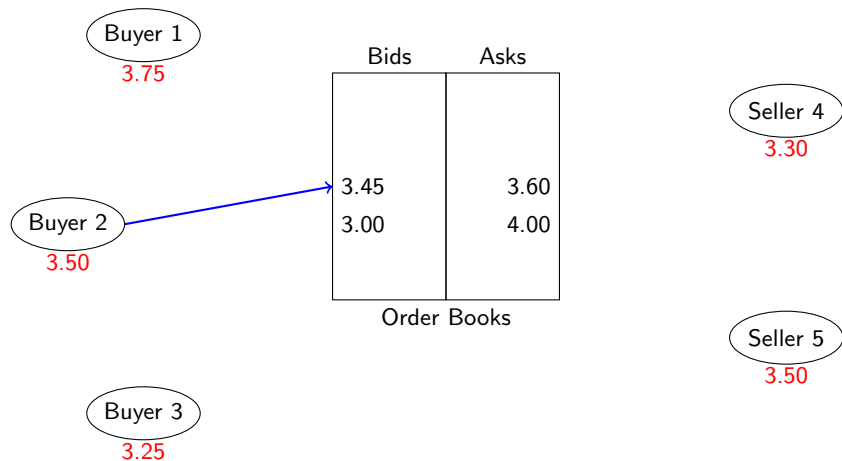
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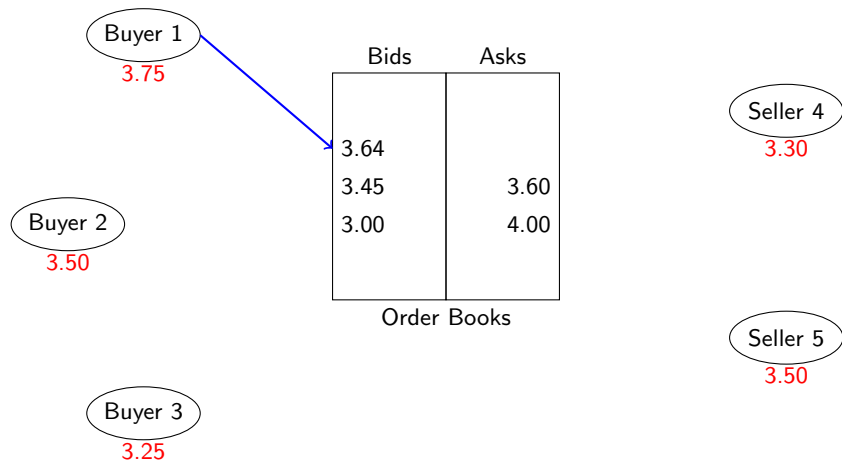
Example



Example



Example



Example

Buyer 1

3.75

Buyer 2

3.50

Buyer 3

3.25

Bids	Asks
3.64	
3.45	3.60
3.00	4.00

Order Books

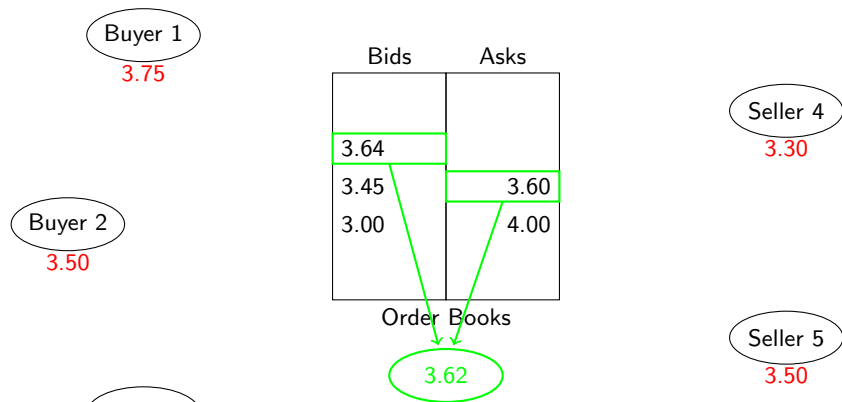
Seller 4

3.30

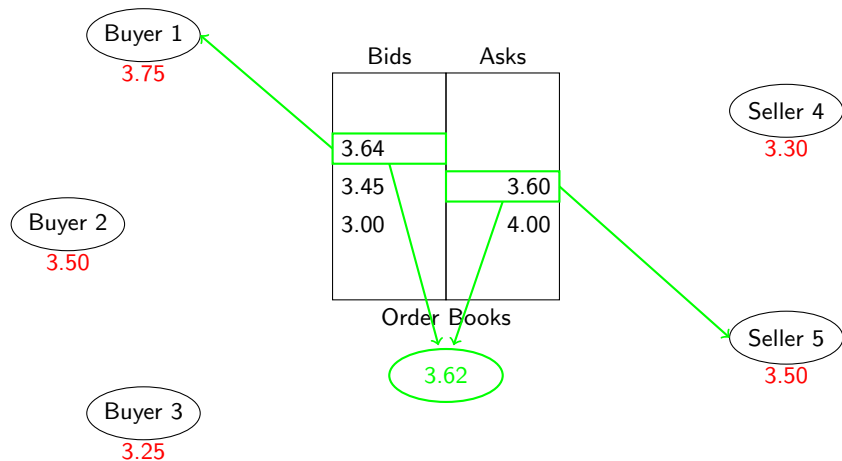
Seller 5

3.50

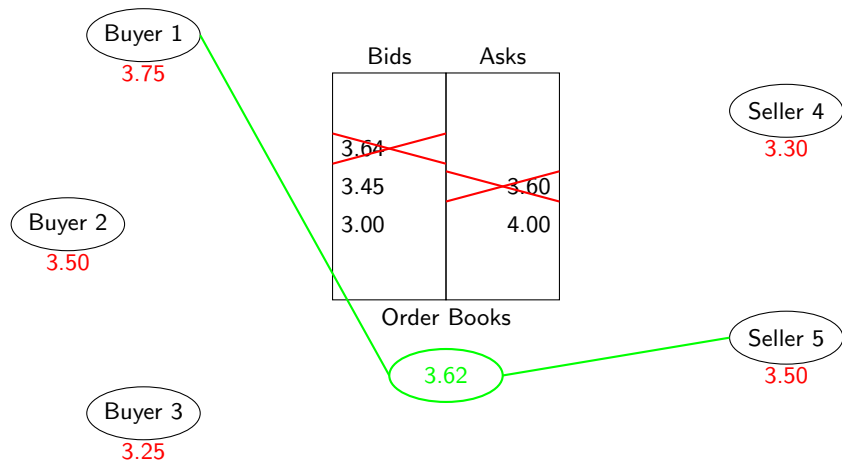
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Execution Uncertainty

- Buyer 1 and Seller 5 transact at 3.62
- Buyer 1 has limit price 3.75

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- So, Buyer 1 expects to make a loss!
- Hence, the CDA fails when we introduce execution uncertainty

Trust-Based CDA

Commitment in the CDA

When a bid/ask is submitted in the CDA:

- The submitter commits to a transaction at that price
- With *anyone*

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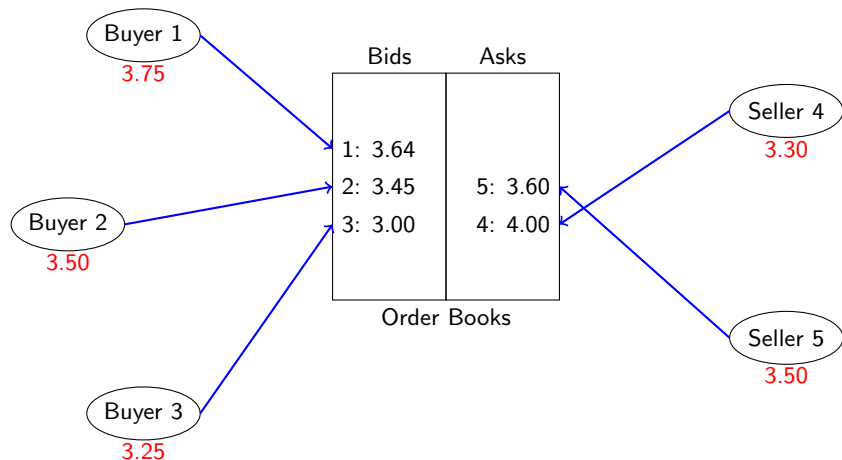
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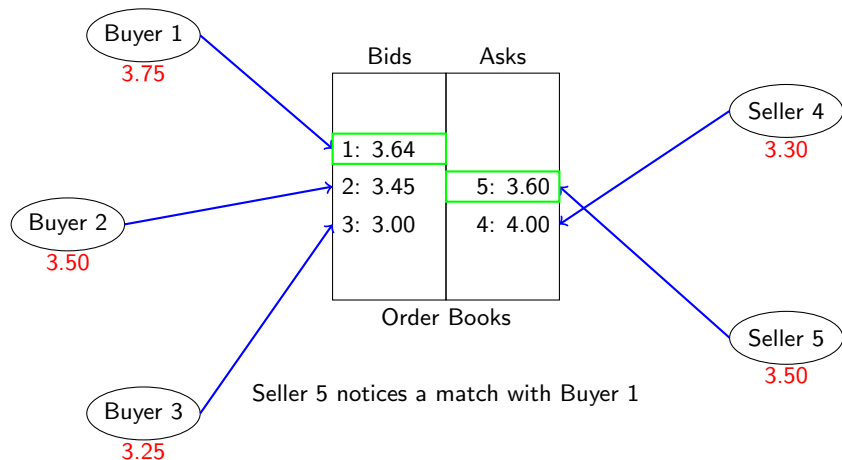
Commitment in the T-CDA

- T-CDA separates commitment from bid/ask
- By adding 'commitment phase': extra step
- Agents can commit to or reject a match
- This allows them to use their *trust* information

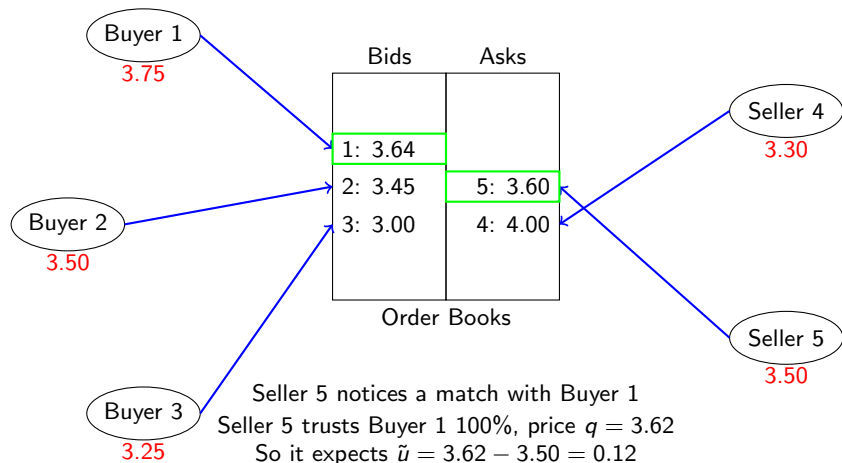
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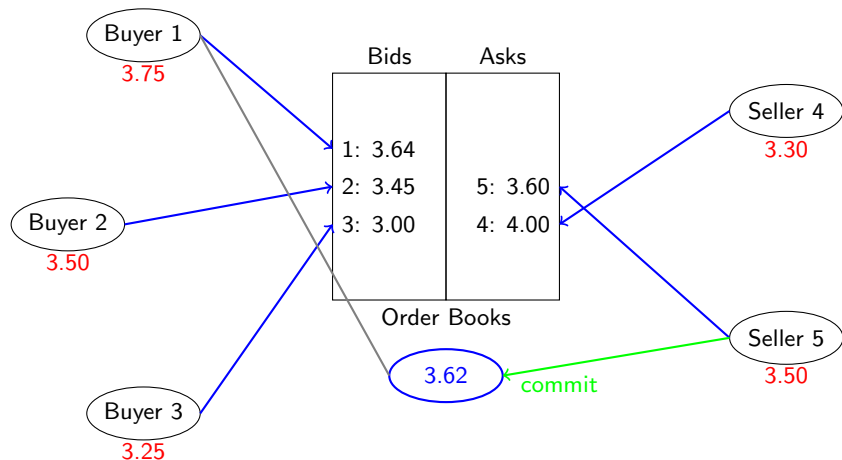
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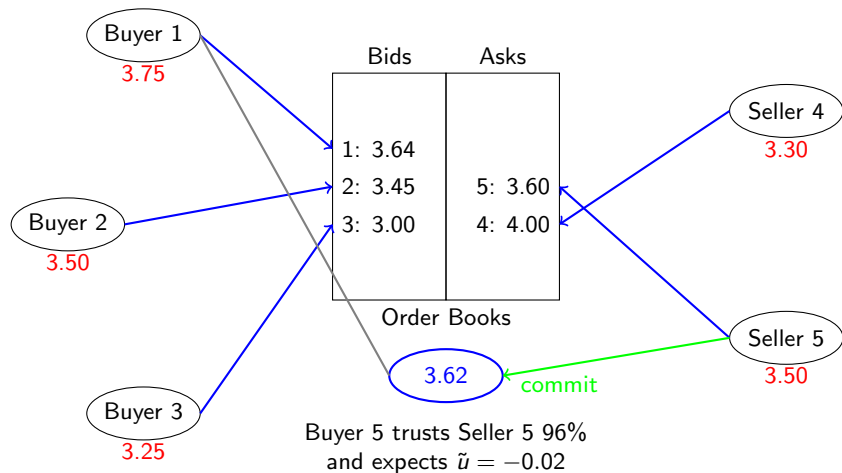
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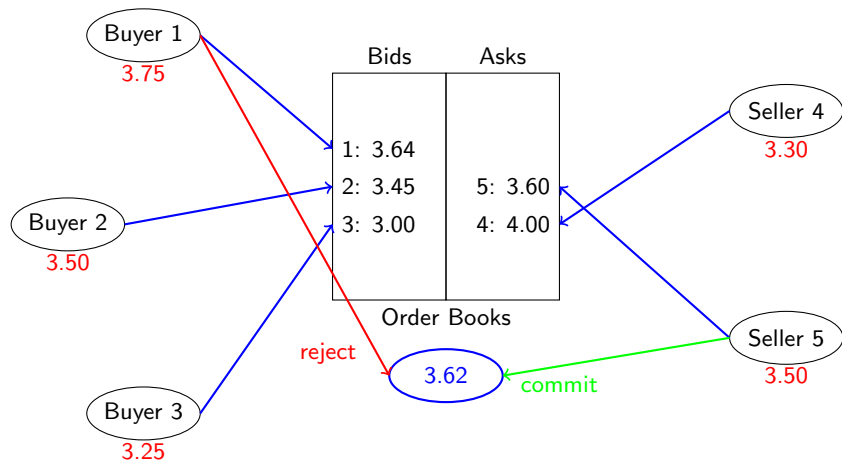
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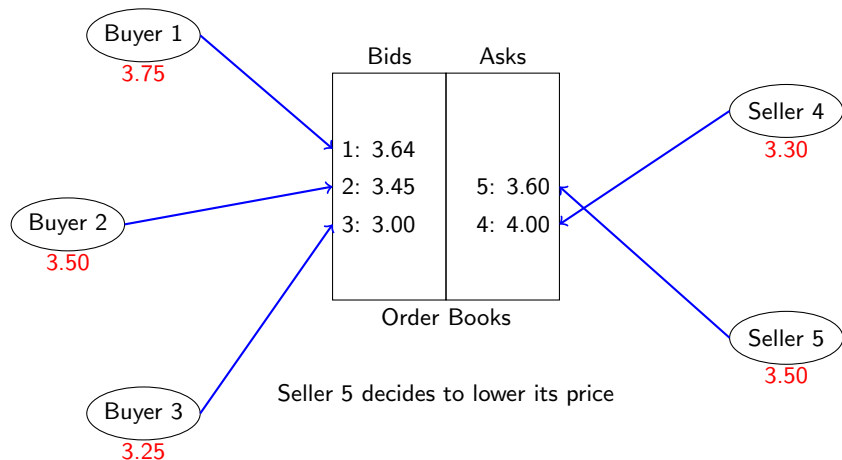
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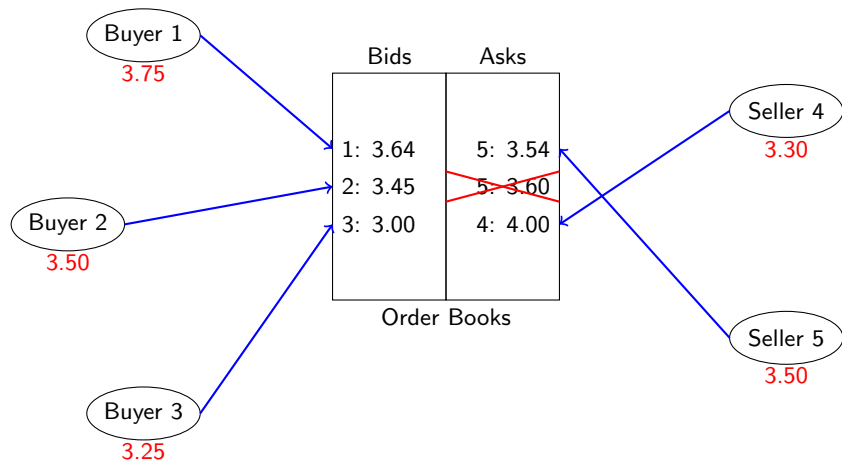
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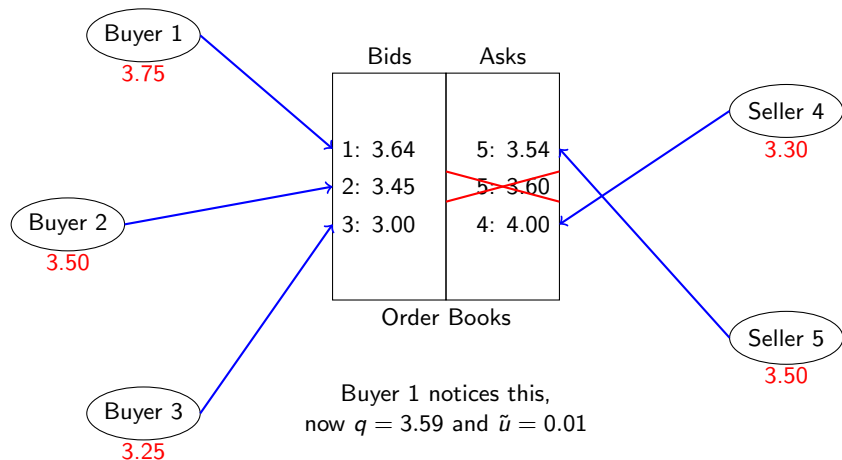
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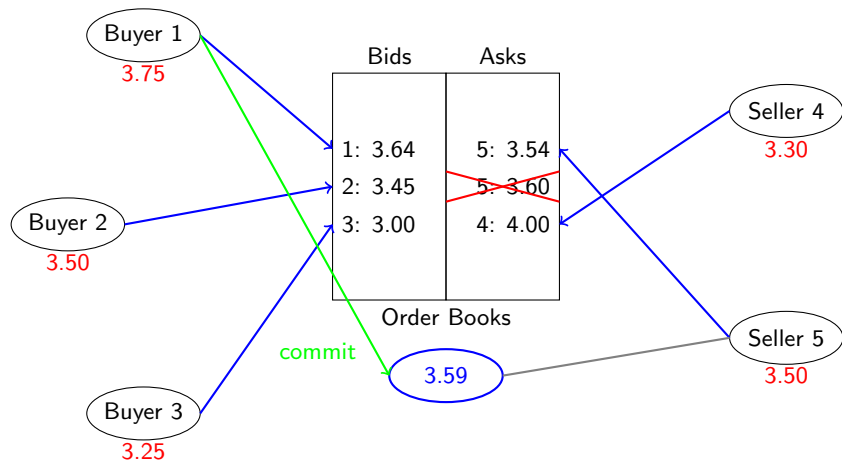
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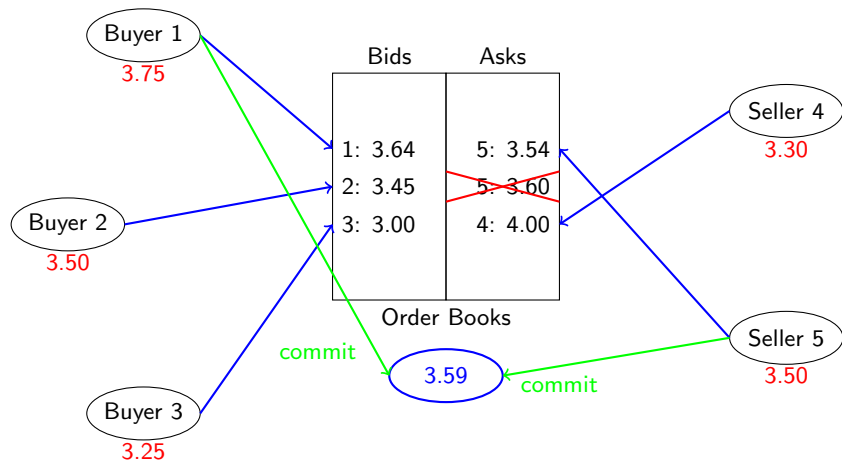
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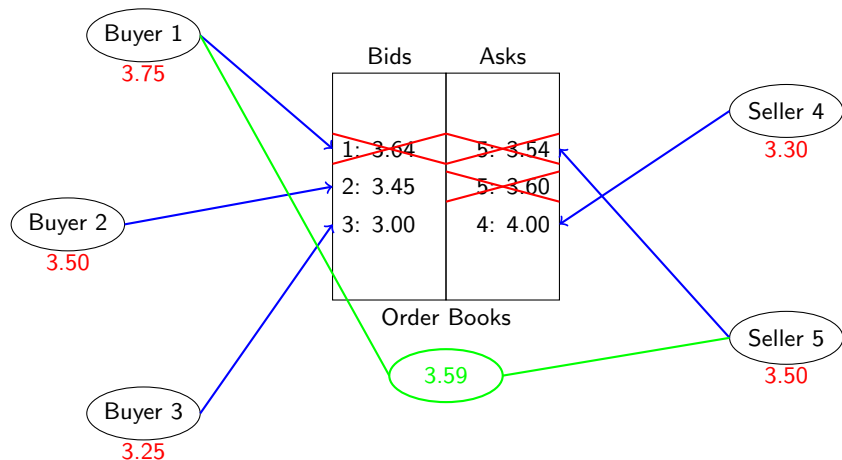
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Evaluation: Strategies (1 of 2)

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- Sellers: uniform from $[\ell, q_{\max}]$

Evaluation: Strategies (2 of 2)

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- All market events are information about market conditions
- Based on bids/asks/commits, an agent updates *profit margin*
- Profit margin and limit price determine bid/ask price
- Prices converge to an *equilibrium price*
- However, in the T-CDA, this does not always exist

Evaluation: Overview

- Assume all traders are given their trust information (fixed values)
- ZI traders:
 - T-CDA vs CDA
 - T-CDA with imperfect information
 - T-CDA vs optimal centralised solution
- T-ZIP traders:
 - Positive: near 100% efficient when $Var(pos) = 0$
 - Negative: fails when $Var(pos) > 0$

T-CDA vs CDA

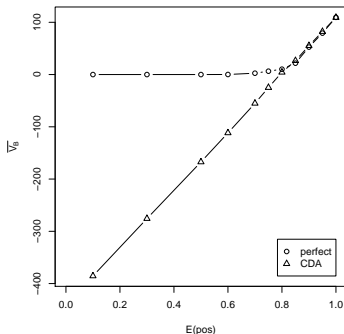


Figure: Utility of the T-CDA (given perfect information) versus CDA, with ZI agents

T-CDA vs CDA

- Traders in the CDA make a loss
- The T-CDA prevents this
- And is always at least as good as the CDA

Imperfect information

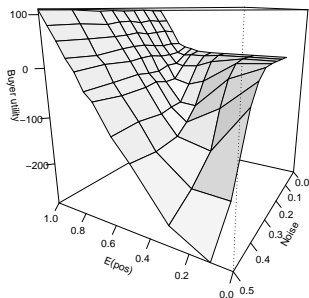


Figure: Performance degrades when $E(pos)$ is lowered and when more noise is added to trust values.

Imperfect information

- Performance degrades linearly when information becomes less accurate

T-CDA vs Optimum

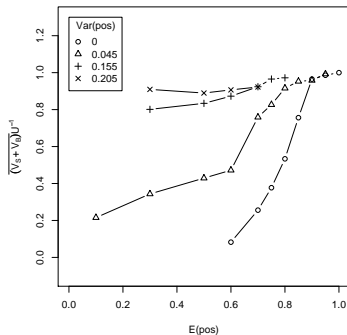


Figure: The normalised utility, or *efficiency* derived by the T-CDA.

T-CDA vs Optimum

- T-CDA with ZI-traders is inefficient for low POS

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T-CDA vs Optimum

- T-CDA with ZI-traders is inefficient for low POS
- ZI bidding range too wide when POS is low
- Hypothesis: this causes inefficiency

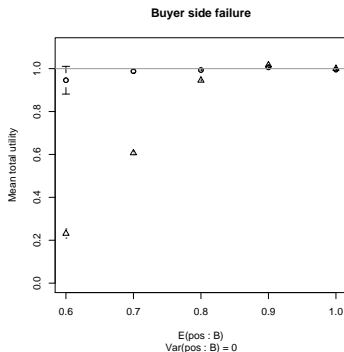
T-ZIP: $Var(pos) = 0$ 

Figure: Buyer side failure, with $Var(pos) = 0$, comparing T-ZIP (circles) and ZI (triangles). Error bars indicate the 95% confidence interval.

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- T-ZIP achieves near-100% efficiency when $Var(pos) = 0$

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- T-ZIP achieves near-100% efficiency when $Var(pos) = 0$
- This confirms that ZI inefficiency is not due to T-CDA
- But that ZI bidding is inappropriate

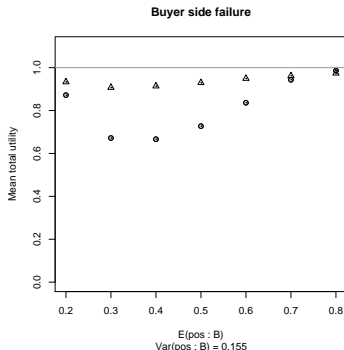
T-ZIP: $Var(pos) > 0$ 

Figure: Buyer side failure, with $Var(pos) = 0.155$, comparing T-ZIP (circles) and ZI (triangles). Error bars indicate the 95% confidence interval.

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- We saw an example where T-ZIP does poorly
- Even ZI does better!

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- T-ZIP converges even when there is no equilibrium price

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- Even ZI does better!
- T-ZIP converges even when there is no equilibrium price
- This causes failure: harmful convergence

T-ZIP: $Var(pos) > 0$

- We saw an example where T-ZIP does poorly
- Even ZI does better!
- T-ZIP converges even when there is no equilibrium price
- This causes failure: harmful convergence
- Mechanics are quite complex, see thesis

Contents

- 1 Introduction
- 2 Mechanism
- 3 Evaluation
- 4 Discussion**

Conclusions

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- No generally successful strategy

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- Develop a model that *learns* trust from market interactions